## **Test and Set Solution**

This algorithm solves the critical section problem for  $n_{\xi}$  processes using a Test and Set instruction (called TaS here). This instruction does the following function atomically:

```
function TaS(var Lock: boolean): boolean;
            begin
                TaS := Lock;
                Lock := true;
            end:
      The solution is:
1
    var waiting: shared array [0..n-1] of boolean;
2
         Lock: shared boolean;
         j: 0..n-1;
3
4
         key: boolean;
5
6
   repeat
                (* process Pi *)
7
         waiting[i] := true;
8
         key := true;
9
         while waiting[i] and key do
10
                key := TaS(Lock);
11
         waiting[i] := false;
         (* critical section goes here *)
12
         j := i + 1 \mod n;
13
         while (j <> i) and not waiting[j] do
14
             j := j + 1 \mod n;
15
         if j = i then
16
             Lock := false
17
18
         else
19
             waiting[j] := false;
    until false;
20
```

- lines 1-2: These are global to all processes, and are all initialized to false.
- **lines 3–4**: These are local to each process i and are uninitialized.
- **lines 6–11**: This is the entry section. Basically, waiting[i] is **true** as long as process i is trying to get into its critical section; if any other process is in that section, then Lock will also be **true**, and process i will loop in lines 9–10. Once process i can go on, it is no longer waiting for permission to enter, and sets waiting[i] to **false** (line 11); it then proceeds into the critical section. Note that Lock is set to **true** by the *TaS* instruction in line 9 that returns **false**.
- lines 13–19: This is the exit section. When process i leaves the critical section, it must choose which other waiting process may enter next. It starts with the process with the next higher index (line 13). It checks each process to see if that process is waiting for access (lines 14–15); if none is, it simply releases the lock (by setting Lock to false; lines 16–17). However, if some other process process j is waiting for entry, process i simply changes waiting[j] to false to allow process j to enter the critical section (lines 18–19).