

Policy Languages

- Express security policies in a precise way
- High-level languages
 - Policy constraints expressed abstractly
- Low-level languages
 - Policy constraints expressed in terms of program options, input, or specific characteristics of entities on system

High-Level Policy Languages

- Constraints expressed independent of enforcement mechanism
- Constraints restrict entities, actions
- Constraints expressed unambiguously
 - Requires a precise language, usually a mathematical, logical, or programming-like language

Example: Web Browser

- Goal: restrict actions of Java programs that are downloaded and executed under control of web browser
- Language specific to Java programs
- Expresses constraints as conditions restricting invocation of entities

Expressing Constraints

- Entities are classes, methods
 - Class: set of objects that an access constraint constrains
 - Method: set of ways an operation can be invoked
- Operations
 - Instantiation: s creates instance of class c : $s \dashv\vdash c$
 - Invocation: $s1$ executes object $s2$: $s1 \dashv\vdash s2$
- Access constraints
 - **deny**(s op x) **when** b
 - While b is true, subject s cannot perform op on (subject or class) x ; empty s means all subjects

DTEL

- Basis: access can be constrained by types
- Combines elements of low-level, high-level policy languages
 - Implementation-level constructs express constraints in terms of language types
 - Constructs do not express arguments or inputs to specific system commands

Example

- Goal: users cannot write to system binaries
- Subjects in administrative domain can
 - User must authenticate to enter that domain
- Subjects belong to domains:
 - *d_user* ordinary users
 - *d_admin* administrative users
 - *d_login* for login
 - *d_daemon* system daemons

Types

- Object types:
 - *t_sysbin* executable system files
 - *t_readable* readable files
 - *t_writable* writable files
 - *t_dte* data used by enforcement mechanisms
 - *t_generic* data generated from user processes
- For example, treat these as partitions
 - In practice, files can be readable and writable; ignore this for the example

Domain Representation

- Sequence
 - First component is list of programs that start in the domain
 - Other components describe rights subject in domain has over objects of a type
(crwd->t_writable)
means subject can create, read, write, and list (search) any object of type t_writable

d_daemon Domain

```
domain d_daemon = (/sbin/init),  
    (crwd->t_writable),  
    (rd->t_generic, t_readable, t_dte),  
    (rxd->t_sysbin),  
    (auto->d_login);
```

- Compromising subject in *d_daemon* domain does not enable attacker to alter system files
 - Subjects here have no write access
- When /sbin/init invokes login program, login program transitions into *d_login* domain

d_admin Domain

```
domain d_admin =
  (/usr/bin/sh, /usr/bin/csh, /usr/bin/ksh),
  (crwxd->t_generic),
  (crwxd->t_readable, t_writable, t_dte,
                                     t_sysbin),
  (sigtstp->d_daemon);
```

- *sigtstp* allows subjects to suspend processes in *d_daemon* domain
- Admin users use a standard command interpreter

d_user Domain

```
domain d_user =  
    (/usr/bin/sh, /usr/bin/csh, /usr/bin/ksh),  
    (crwd->t_generic),  
    (rxd->t_sysbin),  
    (crwd->t_writable),  
    (rd->t_readable, t_dte);
```

- No auto component as no user commands transition out of it
- Users cannot write to system binaries

d_login Domain

```
domain d_login =  
    (/usr/bin/login),  
    (crwd->t_writable),  
    (rd->t_readable, t_generic, t_dte),  
    setauth,  
    (exec->d_user, d_admin);
```

- Cannot execute anything except the transition
 - Only /usr/bin/login in this domain
- *setauth* enables subject to change UID
- *exec* access to *d_user*, *d_admin* domains

Set Up

```
initial_domain = d_daemon;
```

- System starts in *d_daemon* domain

```
assign -r t_generic /;
```

```
assign -r t_writable /usr/var, /dev, /tmp;
```

```
assign -r t_readable /etc;
```

```
assign -r -s dte_t /dte;
```

```
assign -r -s t_sysbin /sbin, /bin,  
                                /usr/bin, /usr/sbin;
```

- These assign initial types to objects
- `-r` recursively assigns type
- `-s` binds type to name of object (delete it, recreate it, still of given type)

Add Log Type

- Goal: users can't modify system logs; only subjects in *d_admin*, new *d_log* domains can

```
type t_readable, t_writable, t_sysbin,  
      t_dte, t_generic, t_log;
```

- New type *t_log*

```
domain d_log =  
  (/usr/sbin/syslogd),  
  (crwd->t_log),  
  (rwd->t_writable),  
  (rd->t_generic, t_readable);
```

- New domain *d_log*

Fix Domain and Set-Up

```
domain d_daemon = (/sbin/init),  
    (crwd->t_writable),  
    (rxd->t_readable),  
    (rd->t_generic, t_dte, t_sysbin),  
    (auto->d_login, d_log);
```

- Subject in *d_daemon* can invoke logging process
 - Can log, but not execute anything

```
assign -r t_log /usr/var/log;  
assign t_writable /usr/var/log/wtmp,  
    /usr/var/log/utmp;
```

- Set type of logs

Low-Level Policy Languages

- Set of inputs or arguments to commands
 - Check or set constraints on system
- Low level of abstraction
 - Need details of system, commands

Example: X Window System

- UNIX X11 Windowing System
- Access to X11 display controlled by list
 - List says what hosts allowed, disallowed access
`xhost +groucho -chico`
- Connections from host groucho allowed
- Connections from host chico not allowed

Example: tripwire

- File scanner that reports changes to file system and file attributes
 - *tw.config* describes what may change
 - `/usr/mab/tripwire +gimnpsu012345678-a`
 - Check everything but time of last access (“-a”)
 - database holds previous values of attributes

Example Database Record

```
/usr/mab/tripwire/README 0 ...../. 100600 45763 1
917 10 33242 .gtPvf .gtPvY .gtPvY 0
.ZD4cc0Wr8i21ZKaI..LUOr3
.0fwo5:hf4e4.8TAqd0V4ubv ?..... ...9b3
1M4GX01xbGIX0oVuGo1h15z3
?:Y9jfa04rdzM1q:egt1APgHk
?.Eb9yo.2zkEh1XKovX1:d0wF0kfAvC
?1M4GX01xbGIX2947jdyrior38h15z3 0
```

- file name, version, bitmask for attributes, mode, inode number, number of links, UID, GID, size, times of creation, last modification, last access, cryptographic checksums

Comments

- System administrators not expected to edit database to set attributes properly
- Checking for changes with tripwire is easy
 - Just run once to create the database, run again to check
- Checking for conformance to policy is harder
 - Need to either edit database file, or (better) set system up to conform to policy, then run tripwire to construct database

Example English Policy

- Computer security policy for academic institution
 - Institution has multiple campuses, administered from central office
 - Each campus has its own administration, and unique aspects and needs
- Authorized Use Policy
- Electronic Mail Policy

Authorized Use Policy

- Intended for one campus (Davis) only
- Goals of campus computing
 - Underlying intent
- Procedural enforcement mechanisms
 - Warnings
 - Denial of computer access
 - Disciplinary action up to and including expulsion
- Written informally, aimed at user community

Electronic Mail Policy

- Systemwide, not just one campus
- Three parts
 - Summary
 - Full policy
 - Interpretation at the campus

Summary

- Warns that electronic mail not private
 - Can be read during normal system administration
 - Can be forged, altered, and forwarded
- Unusual because the policy alerts users to the threats
 - Usually, policies say how to prevent problems, but do not define the threats

Summary

- What users should and should not do
 - Think before you send
 - Be courteous, respectful of others
 - Don't interfere with others' use of email
- Personal use okay, provided overhead minimal
- Who it applies to
 - Problem is UC is quasi-governmental, so is bound by rules that private companies may not be
 - Educational mission also affects application

Full Policy

- Context
 - Does not apply to Dept. of Energy labs run by the university
 - Does not apply to printed copies of email
 - Other policies apply here
- E-mail, infrastructure are university property
 - Principles of academic freedom, freedom of speech apply
 - Access without user's permission requires approval of vice chancellor of campus or vice president of UC
 - If infeasible, must get permission retroactively

Uses of E-mail

- Anonymity allowed
 - Provided it doesn't break laws or other policies
- Can't interfere with others' use of e-mail
 - No spam, letter bombs, e-mailed worms, *etc.*
- Personal e-mail allowed within limits
 - Cannot interfere with university business
 - Such e-mail may be a “university record” subject to disclosure

Security of E-mail

- University can read e-mail
 - Won't go out of its way to do so
 - Allowed for legitimate business purposes
 - Allowed to keep e-mail robust, reliable
- Archiving and retention allowed
 - May be able to recover e-mail from end system (backed up, for example)

Implementation

- Adds campus-specific requirements and procedures
 - Example: “incidental personal use” not allowed if it benefits a non-university organization
 - Allows implementation to take into account differences between campuses, such as self-governance by Academic Senate
- Procedures for inspecting, monitoring, disclosing e-mail contents
- Backups

Confidentiality Policy

- Goal: prevent the unauthorized disclosure of information
 - Deals with information flow
 - Integrity incidental
- Multi-level security models are best-known examples
 - Bell-LaPadula Model basis for many, or most, of these

Bell-LaPadula Model, Step 1

- Security levels arranged in linear ordering
 - Top Secret: highest
 - Secret
 - Confidential
 - Unclassified: lowest
- Levels consist of *security clearance* $L(s)$
 - Objects have *security classification* $L(o)$

Example

<i>security level</i>	<i>subject</i>	<i>object</i>
Top Secret	Tamara	Personnel Files
Secret	Samuel	E-Mail Files
Confidential	Claire	Activity Logs
Unclassified	Ulaley	Telephone Lists

- Tamara can read all files
- Claire cannot read Personnel or E-Mail Files
- Ulaley can only read Telephone Lists

Reading Information

- Information flows *up*, not *down*
 - “Reads up” disallowed, “reads down” allowed
- Simple Security Condition (Step 1)
 - Subject s can read object o iff, $L(o) \leq L(s)$ and s has permission to read o
 - Note: combines mandatory control (relationship of security levels) and discretionary control (the required permission)
 - Sometimes called “no reads up” rule

Writing Information

- Information flows up, not down
 - “Writes up” allowed, “writes down” disallowed
- *-Property (Step 1)
 - Subject s can write object o iff $L(s) \leq L(o)$ and s has permission to write o
 - Note: combines mandatory control (relationship of security levels) and discretionary control (the required permission)
 - Sometimes called “no writes down” rule