

Outline for January 12, 2007

1. Greetings and Felicitations!
2. Take-Grant
 - a. Counterpoint to HRU result
 - b. Symmetry of take and grant rights
 - c. Islands (maximal subject-only *tg*-connected subgraphs)
 - d. Bridges (as a combination of terminal and initial spans)
3. Sharing
 - a. Definition: $\text{can}\bullet\text{share}(r, \mathbf{x}, \mathbf{y}, G_0)$ true iff there exists a sequence of protection graphs G_0, \dots, G_n such that $G_0 \vdash^* G_n$ using only take, grant, create, remove rules and in G_n , there is an edge from \mathbf{x} to \mathbf{y} labeled r
 - b. Theorem: $\text{can}\bullet\text{share}(r, \mathbf{x}, \mathbf{y}, G_0)$ iff there is an edge from \mathbf{x} to \mathbf{y} labeled r in G_0 , or all of the following hold:
 - i. there is a vertex \mathbf{y}' with an edge from \mathbf{y}' to \mathbf{y} labeled r ;
 - ii. there is a subject \mathbf{y}'' which terminally spans to \mathbf{y}' , or $\mathbf{y}'' = \mathbf{y}'$;
 - iii. there is a subject \mathbf{x}' which initially spans to \mathbf{x} , or $\mathbf{x}' = \mathbf{x}$; and
 - iv. there is a sequence of islands I_1, \dots, I_n connected by bridges for which \mathbf{x}' is in I_1 and \mathbf{y}' is in I_n .
4. Model Interpretation
 - a. ACM very general, broadly applicable; Take-Grant more specific, can model fewer situations
 - b. Theorem: G_0 protection graph with exactly one subject, no edges; R set of rights. Then $G_0 \vdash^* G$ iff G is a finite directed graph containing subjects and objects only, with edges labeled from nonempty subsets of R , and with at least one subject with no incoming edges
 - c. Example: shared buffer managed by trusted third part
5. Stealing
 - a. Definition: $\text{can}\bullet\text{steal}(r, \mathbf{x}, \mathbf{y}, G_0)$ true iff there is no edge from \mathbf{x} to \mathbf{y} labeled r in G_0 , and there exists a sequence of protection graphs G_0, \dots, G_n such that $G_0 \vdash^* G_n$ in which:
 - i. G_n has an edge from \mathbf{x} to \mathbf{y} labeled r
 - ii. There is a sequence of rule applications ρ_1, \dots, ρ_n such that $G_{i-1} \vdash G_i$; and
 - iii. For all vertices \mathbf{v}, \mathbf{w} in G_{i-1} , if there is an edge from \mathbf{v} to \mathbf{y} in G_0 labeled r , then ρ_i is not of the form “ \mathbf{v} grants (r to \mathbf{y}) to \mathbf{w} ”
 - b. Example