## **Outline for January 17, 2007**

- 1. Greetings and Felicitations!
- 2. Stealing
  - a. Definition:  $can \cdot steal(r, \mathbf{x}, \mathbf{y}, G_0)$  true iff there is no edge from  $\mathbf{x}$  to  $\mathbf{y}$  labeled r in  $G_0$ , and there exists a sequence of protection graphs  $G_0, ..., G_n$  such that  $G_0 \vdash^* G_n$  in which:
    - i.  $G_n$  has an edge from **x** to **y** labeled r
    - ii. There is a sequence of rule applications  $\rho_1, ..., \rho_n$  such that  $G_{i-1} \vdash G_i$ ; and
    - iii. For all vertices **v**, **w** in  $G_{i-1}$ , if there is an edge from **v** to **y** in  $G_0$  labeled *r*, then  $\rho_i$  is not of the form "**v** grants (*r* to **y**) to **w**"
  - b. Example
  - c. Theorem:  $can \bullet steal(r, \mathbf{x}, \mathbf{y}, G_0)$  iff all of the following hold:
    - i. there is no edge from  $\mathbf{x}$  to  $\mathbf{y}$  labeled r in  $G_0$ ;
    - ii. there is a subject  $\mathbf{x}'$  which initially spans to  $\mathbf{x}$ , or  $\mathbf{x}' = \mathbf{x}$ ; and
    - iii. there is a vertex s with an edge to y labeled r in  $G_0$  and for which  $can \cdot share(t, \mathbf{x}, \mathbf{s}, G_0)$  holds
- 3. Conspiracy
  - a. Access set
  - b. Deletion set
  - c. Conspiracy graph
  - d. *I*, *T* sets
  - e. Theorem: *can•share*(r,  $\mathbf{x}$ ,  $\mathbf{y}$ ,  $G_0$ ) iff there is a path from some  $h(\mathbf{p}) \in I(\mathbf{x})$  to some  $h(\mathbf{q}) \in T(\mathbf{y})$
- 4. Schematic Protection Model
  - a. Model components
  - b. Link function
  - c. Filter function
  - d. Example: Take-Grant as an instance of SPM
  - e. Create operations and attenuation
  - f. Flow functions, maximal state
  - g. Safety analysis