ECS 235B Module 34
Introduction to Noninterference

Interference

- Think of it as something used in communication
  - Holly/Lara example: Holly interferes with the CPU utilization, and Lara detects it — communication
- Plays role of writing (interfering) and reading (detecting the interference)

Model

- System as state machine
  - Subjects $S = \{ s_i \}$
  - States $\Sigma = \{ \sigma_i \}$
  - Outputs $O = \{ o_i \}$
  - Commands $Z = \{ z_i \}$
  - State transition commands $C = S \times Z$
- Note: no inputs
  - Encode either as selection of commands or in state transition commands

Functions

- State transition function T: $C \times \Sigma \to \Sigma$
  - Describes effect of executing command c in state σ
- Output function P: $C \times \Sigma \to O$
  - Output of machine when executing command c in state σ
- Initial state is σ_0

Example: 2-Bit Machine

- Users Heidi (high), Lucy (low)
- 2 bits of state, *H* (high) and *L* (low)
  - System state is (*H*, *L*) where *H*, *L* are 0, 1
- 2 commands: *xor0*, *xor1* do xor with 0, 1
  - Operations affect *both* state bits regardless of whether Heidi or Lucy issues it

Example: 2-bit Machine

- S = { Heidi, Lucy }
- Σ = { (0,0), (0,1), (1,0), (1,1) }
- C = { *xor0*, *xor1* }

	Input States (*H*, *L*)			
	(0,0)	(0,1)	(1,0)	(1,1)
xor0	(0,0)	(0,1)	(1,0)	(1,1)
xor1	(1,1)	(1,0)	(0,1)	(0,0)

Outputs and States

- T is inductive in first argument, as

  $T(c_0, \sigma_0) = \sigma_1$; $T(c_{i+1}, \sigma_{i+1}) = T(c_{i+1}, \mathrm{T}(c_i, \sigma_i))$
- Let C^* be set of possible sequences of commands in C
- $T^*: C^* \times \Sigma \rightarrow \Sigma$ and

  $c_s = c_0 \ldots c_n \Rightarrow T^*(c_s, \sigma_i) = T(c_n, \ldots, T(c_0, \sigma_i) \ldots)$
- P similar; define $P^*: C^* \times \Sigma \rightarrow O$ similarly

Projection

- $T^*(c_s,\sigma_i)$ sequence of state transitions
- $P^*(c_s,\sigma_i)$ corresponding outputs
- $proj(s, c_s, \sigma_i)$ set of outputs in $P^*(c_s,\sigma_i)$ that subject *s* authorized to see
  - In same order as they occur in $P^*(c_s,\sigma_i)$
  - Projection of outputs for *s*
- Intuition: list of outputs after removing outputs that *s* cannot see

Purge

- $G \subseteq S$, G a group of subjects
- $A \subseteq Z$, A a set of commands
- $\pi_G(c_s)$ subsequence of c_s with all elements (s,z), $s \in G$ deleted
- $\pi_A(c_s)$ subsequence of c_s with all elements (s,z), $z \in A$ deleted
- $\pi_{G,A}(c_s)$ subsequence of c_s with all elements (s,z), $s \in G$ and $z \in A$ deleted

Example: 2-bit Machine

- Let σ_0 = (0,1)
- 3 commands applied:
  - Heidi applies *xor0*
  - Lucy applies *xor1*
  - Heidi applies *xor1*
- c_s = ((Heidi, *xor0*), (Lucy, *xor1*), (Heidi, *xor1*))
- Output is 011001
  - Shorthand for sequence (0,1) (1,0) (0,1)

Example

- *proj*(Heidi, c_s, σ_0) = 011001
- *proj*(Lucy, c_s, σ_0) = 101
- $\pi_{\text{Lucy}}(c_s)$ = (Heidi, *xor0*), (Heidi, *xor1*)
- $\pi_{\text{Lucy},xor1}(c_s)$ = (Heidi, *xor0*), (Heidi, *xor1*)
- $\pi_{\text{Heidi}}(c_s)$ = (Lucy, *xor1*)
- $\pi_{\text{Lucy},xor0}(c_s)$ = (Heidi, *xor0*), (Lucy, *xor1*), (Heidi, *xor1*)
- $\pi_{\text{Heidi},xor0}(c_s) = \pi_{xor0}(c_s)$ = (Lucy, *xor1*), (Heidi, *xor1*)
- $\pi_{\text{Heidi},xor1}(c_s)$ = (Heidi, *xor0*), (Lucy, *xor1*)
- $\pi_{xor1}(c_s)$ = (Heidi, *xor0*)

Noninterference

- Intuition: If set of outputs Lucy can see corresponds to set of inputs she can see, there is no interference
- Formally: G, $G' \subseteq$ S, $G \neq G'$; $A \subseteq Z$; users in G executing commands in A are *noninterfering* with users in G' iff for all $c_s \in C^*$, and for all $s \in G'$,

$$proj(s, c_s, \sigma_i) = proj(s, \pi_{G,A}(c_s), \sigma_i)$$

  - Written A,G :| G'

Example: 2-Bit Machine

- Let c_s = ((Heidi, *xor0*), (Lucy, *xor1*), (Heidi, *xor1*)) and $\sigma_0 = (0, 1)$
  - As before
- Take G = { Heidi }, G' = { Lucy }, $A = \varnothing$
- $\pi_{\text{Heidi}}(c_s)$ = (Lucy, *xor1*)
  - So $proj(\text{Lucy}, \pi_{\text{Heidi}}(c_s), \sigma_0) = 0$
- $proj(\text{Lucy}, c_s, \sigma_0) = 101$
- So { Heidi } :| { Lucy } is false
  - Makes sense; commands issued to change H bit also affect L bit

Example

- Same as before, but Heidi's commands affect *H* bit only, Lucy's the *L* bit only
- Output is $0_H 0_L 1_H$
- $\pi_{\text{Heidi}}(c_s)$ = (Lucy, *xor1*)
  - So $proj(\text{Lucy}, \pi_{\text{Heidi}}(c_s), \sigma_0) = 0$
- $proj(\text{Lucy}, c_s, \sigma_0) = 0$
- So { Heidi } :| { Lucy } is true
  - Makes sense; commands issued to change *H* bit now do not affect *L* bit

Quiz

Which of the following best describes noninterference *informally*?

1. Someone operating at LOW cannot see HIGH outputs
2. Someone operating at HIGH cannot see LOW outputs
3. When the LOW inputs are the same, the LOW outputs are the same regardless of the HIGH inputs and outputs
4. When the LOW inputs are the same, different HIGH inputs and outputs will affect the LOW outputs