## ECS 235B Module 44 Policy Composition I

## Policy Composition I

- Assumed: Output function of input
- Means deterministic (else not function)
- Means uninterruptability (differences in timings can cause differences in states, hence in outputs)
- This result for deterministic, noninterference-secure systems


## Compose Systems

- Louie, Dewey LOW
- Hughie HIGH
- $b_{L}$ output buffer
- Anyone can read it
- $b_{H}$ input buffer
- From HIGH source
- Hughie reads from:

- $b_{L H}$ (Louie writes)
- $b_{\text {LDH }}$ (Louie, Dewey write)
- $b_{\text {DH }}$ (Dewey writes)


## Systems Secure

## - All noninterference-secure

- Hughie has no output
- So inputs don't interfere with it
- Louie, Dewey have no input
- So (nonexistent) inputs don't interfere with outputs



## Security of Composition

- Buffers finite, sends/receives blocking: composition not secure!
- Example: assume $b_{D H}, b_{L H}$ have capacity 1
- Algorithm:

1. Louie (Dewey) sends message to $b_{L H}\left(b_{D H}\right)$

- Fills buffer

2. Louie (Dewey) sends second message to $b_{L H}\left(b_{D H}\right)$
3. Louie (Dewey) sends a 0 (1) to $b_{L}$
4. Louie (Dewey) sends message to $b_{L D H}$

- Signals Hughie that Louie (Dewey) completed a cycle


## Hughie

- Reads bit from $b_{H}$
- If 0 , receive message from $b_{L H}$
- If 1 , receive message from $b_{D H}$
- Receive on $b_{\text {LDH }}$
- To wait for buffer to be filled


## Example

- Hughie reads 0 from $b_{H}$
- Reads message from $b_{L H}$
- Now Louie's second message goes into $b_{L H}$
- Louie completes setp 2 and writes 0 into $b_{L}$
- Dewey blocked at step 1
- Dewey cannot write to $b_{L}$
- Symmetric argument shows that Hughie reading 1 produces a 1 in $b_{L}$
- So, input from $b_{H}$ copied to output $b_{L}$


## Quiz

## True or False: If a machine has no outputs, it is noninterference-secure

