

# Lecture 1, March 30

ECS 235B, Foundations of Computer and Information Security  
Spring Quarter 2026

# Let's Get the Hard Stuff Out of the Way

- Adding the class
  - The department policy controls who gets PTAs
  - Available at <https://cs.ucdavis.edu/graduate/policies>
    - Click on "PTA Process and Expectations"

**I can only issue PTAs in accordance with those directions  
(that is, only when *the department* asks me)**

# General Information

- Who we are
  - Instructor: Matt Bishop, [mabishop@ucdavis.edu](mailto:mabishop@ucdavis.edu)  
Office hours: Tu 12:10pm–1:00pm, ThF 1:30pm–2:20pm
  - TA: Chien Chou, [billin@ucdavis.edu](mailto:billin@ucdavis.edu); office hours: *to be arranged*
- Goals:
  - Learn about the reference monitor and high assurance systems;
  - Learn about the access control matrix model and its variants, and how it is used to analyze the security of classes of systems;
  - Learn about the mathematics underlying security policies and their composition;
  - Learn about the confinement problem and information flow; and
  - Explore other topics of interest.

# General Information

- Prerequisites
  - ECS 235A (Computer and Information Security), or knowledge of security commensurate with it
  - Recommended: ECS 150 (Operating Systems); ECS 120 (Theory of Computation)
- Texts:
  - M. Bishop, *Computer Security: Art and Science*, 2<sup>nd</sup> Edition, Addison-Wesley Professional, Boston, MA (2018). ISBN: 978-0-321-71233-2.
  - *Recommended*: R. Anderson, *Security Engineering: A Guide to Building Dependable Distributed Systems*, 3<sup>rd</sup> Edition, John Wiley & Sons, Inc., New York, NY (2020). ISBN: 978-1-119-64281-7.

# General Information

- Class Web Site
  - It's on Canvas; log in there, and go to ECS 235B
  - If you don't have access and need it, please email me and I will add you
    - **This *does not* mean you will be admitted to the course!**
  - Backup web site: <http://nob.cs.ucdavis.edu/classes/ecs235b-2026-02>
- Grading
  - Homework is 50% total, and I expect 4 or 5 equally weighted assignments
  - Project is 50%, weighted as described in the handout

# General Information

- General UC Davis resources
  - Cover health and wellness, virtual classroom fatigue, etc.
  - You can also come to me and I will either help or suggest where you can get help
  - Web site: <https://ebeler.faculty.ucdavis.edu/resources/faq-student-resources/>
- Academic integrity
  - ***Paramount!!!*** And plagiarism is a violation of academic integrity
  - Follow the UC Davis Code of Academic Conduct
    - Available at <https://ossja.ucdavis.edu/code-academic-conduct>
  - I report cheating to the Office of Student Support and Judicial affairs; they handle it appropriately

# Homework

- Due at 11:55pm on the due date, unless stated otherwise
- Submit it on Canvas
  - We will give you comments and grades on Canvas too
  - None of this will be on the secondary web site; only on Canvas
- Don't be wishy-washy in your answer
- Write it in good English; be clear and concise
  - An example of a good answer is on the web site
- Submit PDF or text files *only*; do *not* submit other formats!
  - Other formats often don't render well across all systems

# Homework

- Late policy: I will accept homework up to 5 week days late
  - You will lose 20% per weekday late
  - Example: your homework is due Monday and you turn it in Wednesday; that is 2 days late, so your actual score will be 60% of what it would have been had you turned it in on time
  - Extensions granted on request, for good reason
- Extra credit
  - It is *not* added into your homework score
  - It is used *only* if your grade is on a borderline (for example, right between an A- and an A); it helps determine which one you get

# Term Project

- Goal: get you to explore an area of security in depth
- It can be:
  - Detailed survey or research paper
  - Programming project on validating or working with some formalism
  - Implementing a policy model and formally verify your implementation is correct
- Key rule: ***pick something that interests you***
- Team size: no more than 5
  - If for some reason a team wants more than 5 people, talk to me about it
  - The larger the team, the more I will expect from it

# Term Project

Due dates:

- *Project selection*: April 17, 2026; 10% of project score
- *Progress report*: May 11, 2026; 20% of project score
- *Final project*: June 8, 2026, by 5:30pm; 70% of project score

Only one team member should submit these

- *Be sure to name all team members in all submissions!*
- Other team members should submit a short note saying who turned it in (one sentence is fine)
  - Please do not put it in the comments field!

# Class

- Asking questions: speak up!
  - If you have a question, so does at least half the class

*Remember, the only stupid question is the one you want answered and don't ask*

# Basic Components of Security

- Confidentiality
  - Keeping data and resources hidden
- Integrity
  - Data integrity (integrity)
  - Origin integrity (authentication)
- Availability
  - Allowing access to data and resources

# McCumber Cube

## Critical Information Characteristics

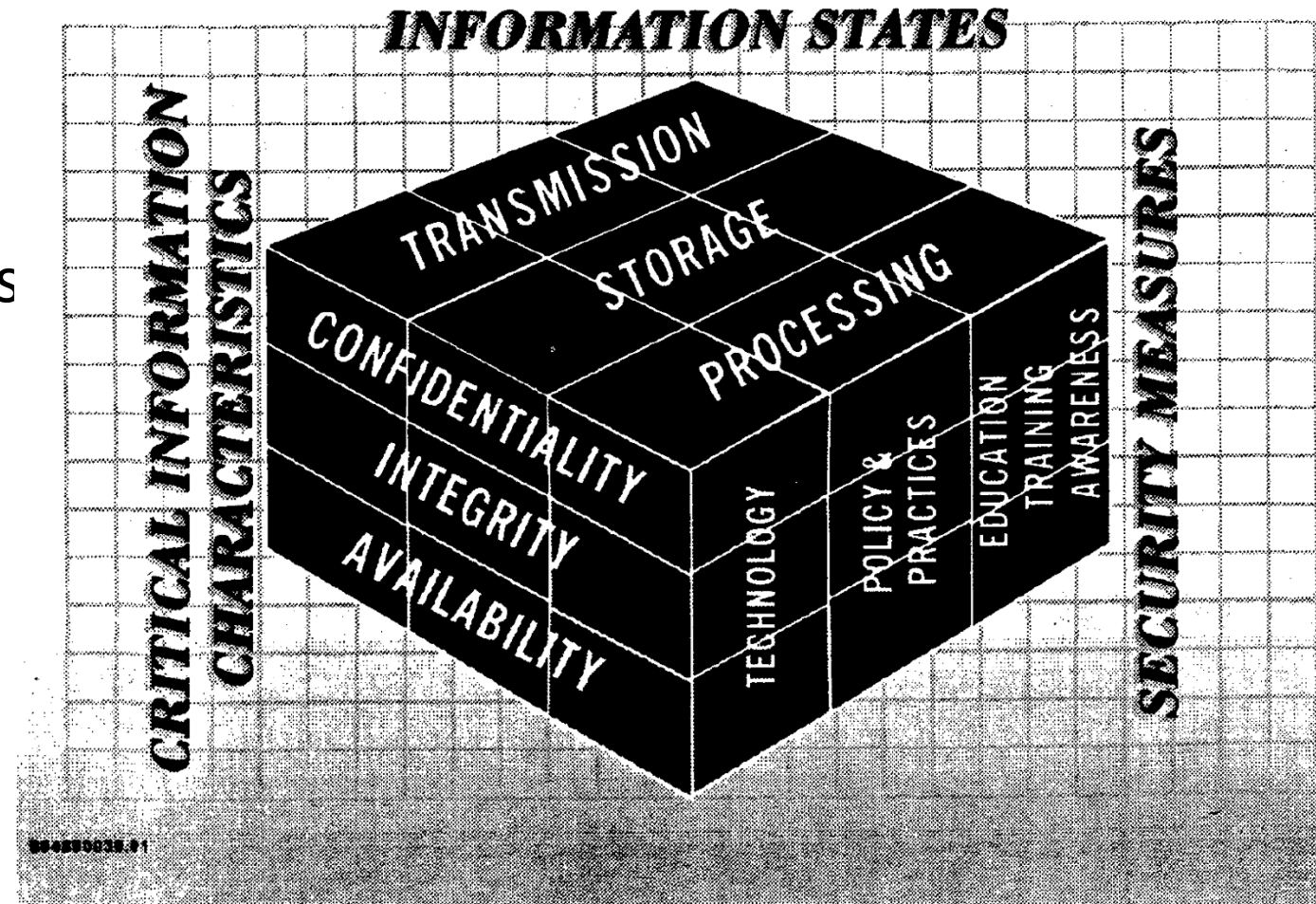
- CIA Triad

## Information States

- Stages in handling information

## Security Measures

- Controls for information access and handling



Picture from J. McCumber, "Information Systems Security: A Comprehensive Model", *Proceedings of the 14th National Computer Security Conference* pp. 328–337 (Oct. 1991); on p. 334.

# Information States

- Storage
  - Where data is kept
  - Examples: disks, USB memory sticks
- Transmission
  - How data moves from one place to another
  - Examples: network connections, pipes
- Processing
  - Computations using the information
  - Examples: computing statistics, drawing pictures

# Security Measures

- Technology
  - Something implemented and used to ensure critical information characteristics maintained through information states
  - Example: encryption, access controls
- Policy and Practice
  - Something which says what information can be accessed, by whom, and how; a procedure to enhance security
  - Example: students may not access one another's homework files
- Education, Training, and Awareness
  - Make people understand security at level appropriate for them
  - Example: cybersecurity training UC Davis folks must take

# Assumptions and Trust

- Underlie *all* aspects of security
- Policies
  - Unambiguously partition system states
  - Correctly capture security requirements
- Mechanisms
  - Assumed to enforce policy
  - Support mechanisms work correctly

# Some Terms: Entities

- **Subject: active entity**
  - Causes information to flow or system state to change
  - Examples: processes, some devices
  - At a higher layer of abstraction: users, other computers
- **Object: passive entity**
  - Contains or receives information
  - Examples: files, some devices
  - At a higher layer of abstraction: file server, network
- **Note: in older papers, objects are entities and so include subjects**
  - This has changed in modern papers

# Reference Monitor

- *Reference monitor* is access control concept of an abstract machine that mediates all accesses to objects by subjects
- *Reference validation mechanism (RVM)* is an implementation of the reference monitor concept.
  - Tamperproof
  - Complete (always invoked and can never be bypassed)
  - Simple (small enough to be subject to analysis and testing, the completeness of which can be assured)
    - Last engenders trust by providing evidence of correctness
- Note: RVM is almost always called a reference monitor too

# Examples

- *Security kernel* combines hardware and software to implement reference monitor
- *Trusted computing base (TCB)* consists of all protection mechanisms within a system responsible for enforcing security policy
  - Includes hardware and software
  - Generalizes notion of security kernel

# Policy and Reference Monitor

- Reference monitor implements a given policy
  - It has a tamperproof authorization database
  - Also maintains an audit trail (record of security-related events) for review
- More on this later; we need some background first

# Access Control Matrix (ACM)

- Basis for representing security policies
- Standard version developed by Harrison, Ruzzo, Ullman to study decidability of security
  - Needed to represent a very general security policy
  - Also a very simple protection system
  - More on this after discussing what we need in an ACM

# Description

objects (entities)

	$O_1$	...	$O_m$	$S_1$	...	$S_n$
$s_1$						
$s_2$						
...						
$s_n$						

subjects

- Subjects  $S = \{ s_1, \dots, s_n \}$
- Objects  $O = \{ o_1, \dots, o_m \}$
- Rights  $R = \{ r_1, \dots, r_k \}$
- Entries  $A[s_i, o_j] \subseteq R$
- $A[s_i, o_j] = \{ r_x, \dots, r_y \}$  means subject  $s_i$  has rights  $r_x, \dots, r_y$  over object  $o_j$

# Example 1

- Processes  $p, q$
- Files  $f, g$
- Rights  $r, w, x, a, o$

	$f$	$g$	$p$	$q$
$p$	$rwo$	$r$	$rwxo$	$w$
$q$	$a$	$ro$	$r$	$rwxo$

# Example 2

- Host names *telegraph*, *nob*, *toadflax*
- Rights *own*, *ftp*, *nfs*, *mail*

	<i>telegraph</i>	<i>nob</i>	<i>toadflax</i>
<i>telegraph</i>	<i>own</i>	<i>ftp</i>	<i>ftp</i>
<i>nob</i>		<i>ftp, mail, nfs, own</i>	<i>ftp, nfs, mail</i>
<i>toadflax</i>		<i>ftp, mail</i>	<i>ftp, mail, nfs, own</i>

# Example 3

- Procedures *inc\_ctr*, *dec\_ctr*, *manage*
- Variable *counter*
- Rights +, −, *call*

	<i>counter</i>	<i>inc_ctr</i>	<i>dec_ctr</i>	<i>manage</i>
<i>inc_ctr</i>	+			
<i>dec_ctr</i>	−			
<i>manager</i>		<i>call</i>	<i>call</i>	<i>call</i>

# UNIX/Linux Access Controls

- Files

- *A* is `~bishop/a.out` (permissions set to 0755, or `rwxr-xr-x`)
- *B* is `/etc/passwd` (permissions set to 0644, or `rw-r--r--`)
- *H* is `/home/bishop` (permissions set to 0711, or `rwX--X--X`)
- *S* is `/bin/su` (permissions set to 4711, or `s--rwx--X--X`)

	<i>A</i>	<i>B</i>	<i>S</i>	<i>H</i>
<i>bishop</i>	<i>rwXO</i>	<i>r</i>	<i>X</i>	<i>rwXO</i>
<i>holly</i>	<i>rx</i>	<i>r</i>	<i>X</i>	<i>X</i>
<i>root</i>	<i>rwX</i>	<i>rWO</i>	<i>rwXO</i>	<i>rwX</i>

# UNIX/Linux Access Controls

- Access control matrices are dynamic:
- After bishop executes `chmod 700 /home/bishop`:
  - Same as `chmod u=rwx,g-rwx,o-rwx /home/bishop`

	<i>A</i>	<i>B</i>	<i>S</i>	<i>H</i>
<i>bishop</i>	<i>rwxo</i>	<i>r</i>	<i>x</i>	<i>rwxo</i>
<i>heidi</i>		<i>rx</i>		
<i>root</i>	<i>rwX</i>	<i>rwo</i>	<i>rwXO</i>	<i>rwX</i>

# Boolean Expression Evaluation

- ACM controls access to database fields
  - Subjects have attributes
  - Verbs define type of access
  - Rules associated with objects, verb pair
- Subject attempts to access object
  - Rule for object, verb evaluated, grants or denies access

# Example

- Subject annie
  - Attributes *role* (artist), *group* (creative)
- Verb paint
  - Default 0 (deny unless explicitly granted)
- Object picture
  - Rule:  
paint: 'artist' in subject.role and  
'creative' in subject.groups and  
time.hour  $\geq 0$  and time.hour  $\leq 4$

# ACM at 3AM and 10AM

At 3AM, time condition met  
ACM is:

... picture ...

...			
annie ...		paint	
...			

At 10AM, time condition not met  
ACM is:

... picture ...

...			
annie ...			
...			

# History

- Problem: what a process has accessed may affect what it can access now
- Example: procedure in a web applet can access other procedures depending on what procedures it has already accessed
  - $S$  set of *static rights* associated with procedure
  - $C$  set of *current rights* associated with each executing process
  - When process calls procedure, rights are  $S \cap C$

# Example Program

```
// This routine has no filesystem access rights  
// beyond those in a limited, temporary area
```

```
procedure helper_proc()  
    return sys_kernel_file
```

```
// But this has the right to delete files
```

```
program main()  
    sys_load_file(helper_proc)  
    tmp_file = helper_proc()  
    sys_delete_file(tmp_file)
```

- *sys\_kernel\_file* contains system kernel
- *tmp\_file* is in limited area that *helper\_proc()* can access it

# Before *helper\_proc* Called

- Static rights of program

	<i>sys_kernel_file</i>	<i>tmp_file</i>
<i>main</i>	delete	delete
<i>helper_proc</i>		delete

- When program starts, current rights:

	<i>sys_kernel_file</i>	<i>tmp_file</i>
<i>main</i>	delete	delete
<i>helper_proc</i>		delete
<i>process</i>	delete	delete

# After *helper\_proc* Called

- Process rights are intersection of static, previous “current” rights:

	<i>sys_kernel_file</i>	<i>tmp_file</i>
<i>main</i>	delete	delete
<i>helper_proc</i>		delete
<i>process</i>		delete

# State Transitions

- Change the protection state of system
- $\vdash$  represents transition
  - $X_i \vdash_{\tau} X_{i+1}$ : command  $\tau$  moves system from state  $X_i$  to  $X_{i+1}$
  - $X_i \vdash^* Y$ : a sequence of commands moves system from state  $X_i$  to  $Y$
- Commands often called *transformation procedures*

# Primitive Operations

- **create subject  $s$ ; create object  $o$** 
  - Creates new row, column in ACM; creates new column in ACM
- **destroy subject  $s$ ; destroy object  $o$** 
  - Deletes row, column from ACM; deletes column from ACM
- **enter  $r$  into  $A[s, o]$** 
  - Adds  $r$  rights for subject  $s$  over object  $o$
- **delete  $r$  from  $A[s, o]$** 
  - Removes  $r$  rights from subject  $s$  over object  $o$

# Create Subject

- Precondition:  $s \notin S$
- Primitive command: **create subject  $s$**
- Postconditions:
  - $S' = S \cup \{s\}, O' = O \cup \{s\}$
  - $(\forall y \in O') [A'[s, y] = \emptyset], (\forall x \in S') [A'[x, s] = \emptyset]$
  - $(\forall x \in S)(\forall y \in O) [A'[x, y] = A[x, y]]$

# Create Object

- Precondition:  $o \notin O$
- Primitive command: **create object**  $o$
- Postconditions:
  - $S' = S, O' = O \cup \{o\}$
  - $(\forall x \in S') [A'[x, o] = \emptyset]$
  - $(\forall x \in S)(\forall y \in O) [A'[x, y] = A[x, y]]$

# Add Right

- Precondition:  $s \in S, o \in O$
- Primitive command: **enter  $r$  into  $A[s, o]$**
- Postconditions:
  - $S' = S, O' = O$
  - $A'[s, o] = A[s, o] \cup \{r\}$
  - $(\forall x \in S')(\forall y \in O' - \{o\}) [A'[x, y] = A[x, y]]$
  - $(\forall x \in S' - \{s\})(\forall y \in O') [A'[x, y] = A[x, y]]$

# Delete Right

- Precondition:  $s \in S, o \in O$
- Primitive command: **delete  $r$  from  $A[s, o]$**
- Postconditions:
  - $S' = S, O' = O$
  - $A'[s, o] = A[s, o] - \{r\}$
  - $(\forall x \in S')(\forall y \in O' - \{o\}) [A'[x, y] = A[x, y]]$
  - $(\forall x \in S' - \{s\})(\forall y \in O') [A'[x, y] = A[x, y]]$

# Destroy Subject

- Precondition:  $s \in S$
- Primitive command: **destroy subject  $s$**
- Postconditions:
  - $S' = S - \{s\}, O' = O - \{s\}$
  - $(\forall y \in O') [A'[s, y] = \emptyset], (\forall x \in S') [A'[x, s] = \emptyset]$
  - $(\forall x \in S')(\forall y \in O') [A'[x, y] = A[x, y]]$

# Destroy Object

- Precondition:  $o \in O$
- Primitive command: **destroy object  $o$**
- Postconditions:
  - $S' = S, O' = O - \{ o \}$
  - $(\forall x \in S') [A'[x, o] = \emptyset]$
  - $(\forall x \in S')(\forall y \in O') [A'[x, y] = A[x, y]]$

# Creating File

- Process  $p$  creates file  $f$  with  $r$  and  $w$  permission

```
command create•file( $p$ ,  $f$ )  
    create object  $f$ ;  
    enter own into  $A[p, f]$ ;  
    enter  $r$  into  $A[p, f]$ ;  
    enter  $w$  into  $A[p, f]$ ;  
end
```

# Mono-Operational Commands

- Make process  $p$  the owner of file  $g$

**command** *make* • *owner*( $p$ ,  $g$ )

**enter** *own* **into**  $A[p, g]$  ;

**end**

- Mono-operational command
  - Single primitive operation in this command

# Conditional Commands

- Let  $p$  give  $q$   $r$  rights over  $f$ , if  $p$  owns  $f$

**command**  $grant \cdot read \cdot file \cdot 1(p, f, q)$

**if**  $own$  **in**  $A[p, f]$

**then**

**enter**  $r$  **into**  $A[q, f];$

**end**

- Monoconditional command
  - Single condition in this command

# Conditions with and

- Let  $p$  give  $q$   $r$  and  $w$  rights over  $f$ , if  $p$  owns  $f$  and  $p$  has  $c$  rights over  $q$

```
command grant•read•file•2(p, f, q)  
    if own in  $A[p, f]$  and  $c$  in  $A[p, q]$   
    then  
        enter  $r$  into  $A[q, f]$ ;  
        enter  $w$  into  $A[q, f]$ ;  
end
```

- Biconditional command
  - Two conditions in this command

# There Is No or

- Let  $p$  give  $q$   $r$  and  $w$  rights over  $f$ , if  $p$  owns  $f$  or  $p$  has  $c$  rights over  $q$

```
command grant.read.file.3( $p, f, q$ )
  if own in  $A[p, f]$ 
  then
    enter  $r$  into  $A[q, f]$ ;
    enter  $w$  into  $A[q, f]$ ;
  end
command grant.read.file.4( $p, f, q$ )
  if  $c$  in  $A[p, q]$ 
  then
    enter  $r$  into  $A[q, f]$ ;
    enter  $w$  into  $A[q, f]$ ;
  end
grant.read.file.3( $p, f, q$ );
grant.read.file.4( $p, f, q$ )
```

# General Form

```
command name of command(parameters)  
    if condition [ and condition [ and condition ... ] ]  
    then  
        list of commands, primitive operations  
end
```

- Only one **if**; no nested **ifs** either
- It must come *before* any primitive operations or subcommands
- No commands may follow the **if**
- There is no **else**

# Copy Flag and Right

- Allows possessor to give rights to another
- Can be a right itself, in which case any other right of the subject can be copied
- Often attached to a right (called a *flag*), so only applies to that right
  - $r$  is read right that cannot be copied
  - $r^c$  is read right that can be copied
- Is copy flag copied when giving  $r$  rights?
  - Depends on model, instantiation of model

# Own Right

- Distinguished right allowing possessor to change entries in ACM column
  - So owner of object can add, delete rights for others
  - May depend on what system allows
    - Can't give rights to specific (set of) users
    - Can't pass copy flag to specific (set of) users

# Principle of Attenuation of Privilege

- Subjects can't increase their rights, or give rights they do not possess
  - Restricts addition of rights within a system
  - Usually *ignored* for owner
    - Why? Owner gives itself rights, gives them to others, deletes the rights it gave itself